

Complementary Slackness Conditions

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1 Complementary Slackness

In the previous notes we established weak duality and used a worked example to motivate strong duality. In this section we describe *complementary slackness*, which gives a precise relationship between optimal primal solutions and optimal dual solutions. Complementary slackness provides an optimality test and explains how to read dual information directly from a final primal tableau (and vice versa).

1.1 Primal and dual slacks

Consider the primal–dual pair in standard form

$$\begin{aligned} \text{maximize} \quad & z = \sum_{j=1}^n c_j x_j \\ \text{subject to} \quad & \sum_{j=1}^n a_{i,j} x_j \leq b_i, \quad 1 \leq i \leq m, \\ & x_j \geq 0, \quad 1 \leq j \leq n \end{aligned} \tag{1}$$

and

$$\begin{aligned} \text{minimize} \quad & w = \sum_{i=1}^m b_i y_i \\ \text{subject to} \quad & \sum_{i=1}^m a_{i,j} y_i \geq c_j, \quad 1 \leq j \leq n, \\ & y_i \geq 0, \quad 1 \leq i \leq m \end{aligned} \tag{2}$$

Definition 1. Let \mathbf{x} be primal feasible and let \mathbf{y} be dual feasible.

(a) The *primal slack vector* is $\mathbf{s} = \mathbf{b} - A\mathbf{x} \geq 0$.

(b) The *dual slack vector* is $\mathbf{t} = A^T \mathbf{y} - \mathbf{c} \geq 0$.

Thus, the i th primal constraint is *tight* if and only if $s_i = 0$, and the j th dual constraint is tight if and only if $t_j = 0$.

1.2 Complementary slackness

The following theorem provide a characterization of primal and dual optimal solutions.

Theorem 2 (Complementary Slackness). *Let \mathbf{x} be primal feasible and let \mathbf{y} be dual feasible. Then \mathbf{x} and \mathbf{y} are optimal solutions to (1) and (2) if and only if*

$$y_i \left(b_i - \sum_{j=1}^n a_{i,j} x_j \right) = 0, \quad 1 \leq i \leq m, \quad (3)$$

$$x_j \left(\sum_{i=1}^m a_{i,j} y_i - c_j \right) = 0, \quad 1 \leq j \leq n. \quad (4)$$

Proof. Let \mathbf{x} be primal feasible and \mathbf{y} dual feasible. Since $\mathbf{s} = \mathbf{b} - A\mathbf{x} \geq 0$, $\mathbf{t} = A^T\mathbf{y} - \mathbf{c} \geq 0$, and $\mathbf{x}, \mathbf{y} \geq 0$, we have

$$\mathbf{y}^T \mathbf{s} \geq 0 \quad \text{and} \quad \mathbf{x}^T \mathbf{t} \geq 0.$$

Now expand the primal–dual gap:

$$\begin{aligned} \mathbf{b}^T \mathbf{y} - \mathbf{c}^T \mathbf{x} &= \mathbf{y}^T \mathbf{b} - \mathbf{y}^T A\mathbf{x} + \mathbf{y}^T A\mathbf{x} - \mathbf{c}^T \mathbf{x} \\ &= \mathbf{y}^T (\mathbf{b} - A\mathbf{x}) + \mathbf{x}^T (A^T \mathbf{y} - \mathbf{c}) \\ &= \mathbf{y}^T \mathbf{s} + \mathbf{x}^T \mathbf{t} \\ &= \sum_{i=1}^m y_i s_i + \sum_{j=1}^n x_j t_j. \end{aligned}$$

Each term in the final expression is nonnegative.

(\Rightarrow) If \mathbf{x} and \mathbf{y} are optimal, then by strong duality $\mathbf{c}^T \mathbf{x} = \mathbf{b}^T \mathbf{y}$, and hence the sum of nonnegative terms above is 0. Therefore, $y_i s_i = 0$ for each i and $x_j t_j = 0$ for each j , which is exactly (3)–(4).

(\Leftarrow) If (3)–(4) hold, then $\mathbf{y}^T \mathbf{s} = 0$ and $\mathbf{x}^T \mathbf{t} = 0$, so $\mathbf{b}^T \mathbf{y} = \mathbf{c}^T \mathbf{x}$. By weak duality this implies \mathbf{x} and \mathbf{y} are optimal. \square

1.3 Interpretation and an optimality test

Complementary slackness can be read as two “either–or” statements.

- (1) For each primal constraint i : either the constraint is slack ($s_i > 0$) and then $y_i = 0$, or $y_i > 0$ and then the constraint must be tight ($s_i = 0$).
- (2) For each primal variable j : either $x_j = 0$, or the corresponding dual constraint must be tight ($t_j = 0$).

Corollary 3 (Certificate of optimality). *If \mathbf{x} is primal feasible, \mathbf{y} is dual feasible, and the complementary slackness conditions (3)–(4) hold, then \mathbf{x} is primal optimal and \mathbf{y} is dual optimal.*

Proof. This is the (\Leftarrow) direction in Theorem 2. \square

1.4 Example (continued)

Consider the primal and dual problems:

$$\begin{aligned} \text{maximize} \quad & z = 22x_1 + 31x_2 + 29x_3 \\ \text{subject to} \quad & x_1 + 4x_2 + 6x_3 \leq 73, \\ & 5x_1 - 2x_2 + 3x_3 \leq 68, \\ & x_1, x_2, x_3 \geq 0. \end{aligned}$$

and

$$\begin{aligned} \text{minimize} \quad & w = 73y_1 + 68y_2 \\ \text{subject to} \quad & y_1 + 5y_2 \geq 22, \\ & 4y_1 - 2y_2 \geq 31, \\ & 6y_1 + 3y_2 \geq 29, \\ & y_1, y_2 \geq 0. \end{aligned}$$

From the final tableaus, the optimal solutions are

$$\mathbf{x}^* = \left(\frac{418}{22}, \frac{297}{22}, 0 \right) = \left(19, \frac{27}{2}, 0 \right), \quad \mathbf{y}^* = \left(\frac{199}{22}, \frac{57}{22} \right),$$

and the optimal value is

$$z^* = w^* = \frac{18403}{22}.$$

Check complementary slackness. First compute the primal slack vector at \mathbf{x}^* :

$$\mathbf{s} = \mathbf{b} - A\mathbf{x}^* = \begin{pmatrix} 73 \\ 68 \end{pmatrix} - \begin{pmatrix} 1 & 4 & 6 \\ 5 & -2 & 3 \end{pmatrix} \begin{pmatrix} 19 \\ 27/2 \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \end{pmatrix}.$$

Thus, $s_1 = s_2 = 0$, so (3) holds.

Next compute the dual slack vector at \mathbf{y}^* :

$$\mathbf{t} = A^T \mathbf{y}^* - \mathbf{c} = \begin{pmatrix} 1 & 5 \\ 4 & -2 \\ 6 & 3 \end{pmatrix} \begin{pmatrix} 199/22 \\ 57/22 \end{pmatrix} - \begin{pmatrix} 22 \\ 31 \\ 29 \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ 727/22 \end{pmatrix}.$$

Therefore, $t_1 = t_2 = 0$ and $t_3 = 727/22 > 0$. Complementary slackness (4) requires $x_3^* t_3 = 0$, which holds since $x_3^* = 0$. Hence, the complementary slackness conditions hold and Corollary 3 certifies that \mathbf{x}^* and \mathbf{y}^* are optimal.

1.5 A tableau perspective

Complementary slackness explains the pattern observed between optimal primal and dual tableaus. At optimality, each basic variable has a positive value given by the right-hand side, each nonbasic variable is 0, and the corresponding dual slack records whether the matching inequality is tight. In particular, the dual slack t_j is positive exactly when the primal variable x_j is forced to be 0, and the primal slack s_i is positive exactly when the dual variable y_i is forced to be 0.